

# L14: Fuzzing

## Software Analysis

IIT Guwahati · August 2026

# Motivating Example

```
int main(int argc, char **argv) {  
    int granted = 0;  
    char pw[8];  
    strcpy(pw, argv[1]);  
    if (strcmp(pw, "pass") == 0) granted = 1;  
    if (granted) printf("Access Granted -> root access\n");  
    else        printf("Access Denied!\n");  
}
```

```
$ ./buggy pass  
Access Granted -> root access
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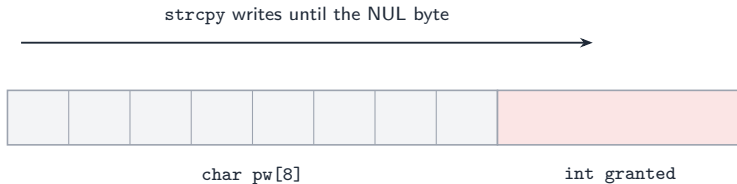
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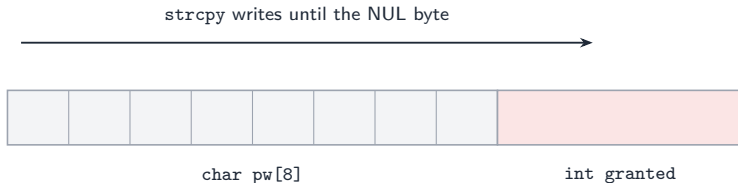
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$ ./buggy abcd  
Access Denied!
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# Why the Long Input Is Accepted



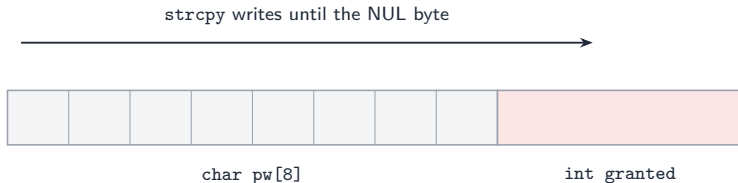
# Why the Long Input Is Accepted



A nine-character argument writes nine bytes — eight characters and the terminating NUL — so the ninth byte lands in `granted` and makes it non-zero:

```
len 8  -> Access Denied!      len 9  -> Access Granted -> root access
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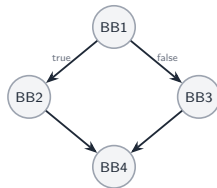
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The program never crashes, so nothing signals a failure: the wrong answer is returned quietly.

# Background: the Control-Flow Graph

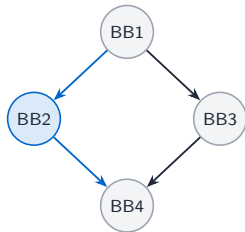
|                                 |                          |
|---------------------------------|--------------------------|
| <code>int check(int x) {</code> | BB1                      |
| <code>    if (x &gt; 3)</code>  | BB1 ends with the branch |
| <code>        x = x - 1;</code> | BB2                      |
| <code>    else</code>           |                          |
| <code>        x = x + 1;</code> | BB3                      |
| <code>    return x;</code>      | BB4                      |
| <code>}</code>                  |                          |



A **basic block** is a straight-line run of instructions with one entry and one exit.

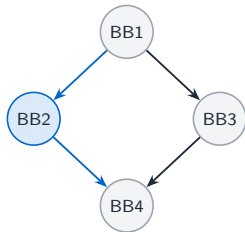


# Execution Paths

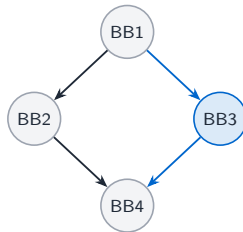


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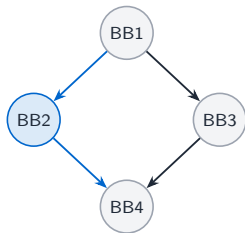


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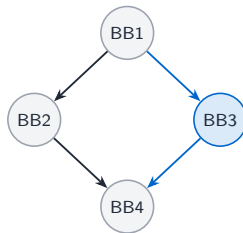


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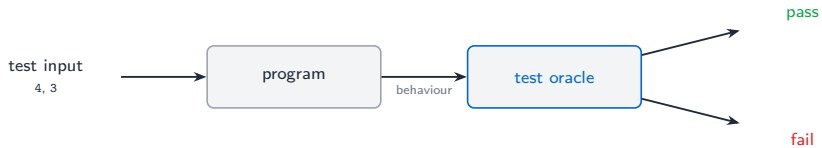
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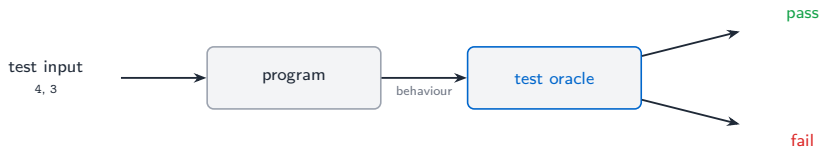
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An **execution path** is one possible flow of control through the program. The two inputs {4,3} execute every block and both branch outcomes: **100% statement and branch coverage** with two test inputs.

# Testcases

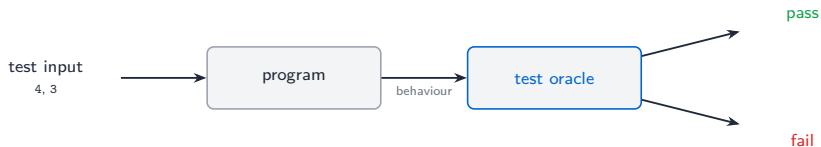


# Testcases



**test case = test input + test oracle**

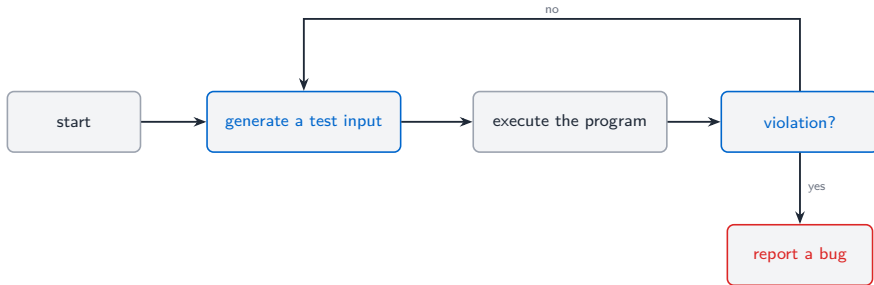
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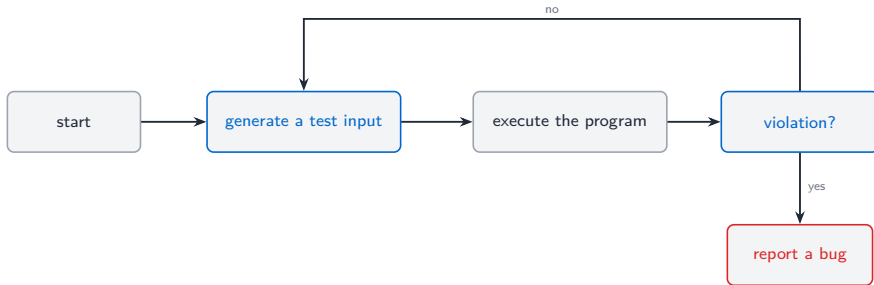
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An input alone only produces behaviour; the oracle is what turns behaviour into a verdict. Automating the two halves raises two different problems.

# The Fuzzing Loop



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The loop is simple; everything of interest is in the two boxes that are not: how an input is **generated**, and what counts as a **violation**.



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*Technical whitepaper for afl-fuzz*

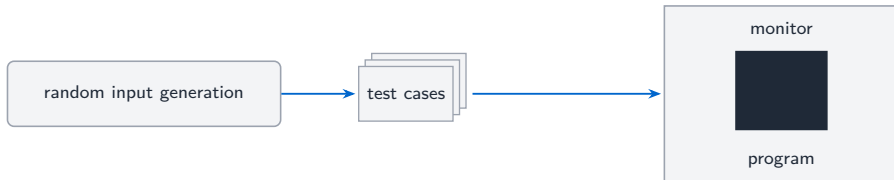
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- **Shallow coverage.** Most generated inputs are rejected early and exercise the same small part of the program.
- **Deep defects are hard to reach.** A defect behind several accepted checks is reached only by an input that passes all of them.
- **Sanity checks block progress.** Magic bytes, checksums, lengths and format tags are exactly the conditions a randomly generated input fails.

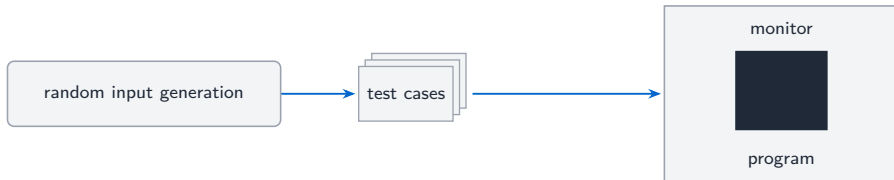
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# Kinds of Fuzzing

# Black-box Fuzzing

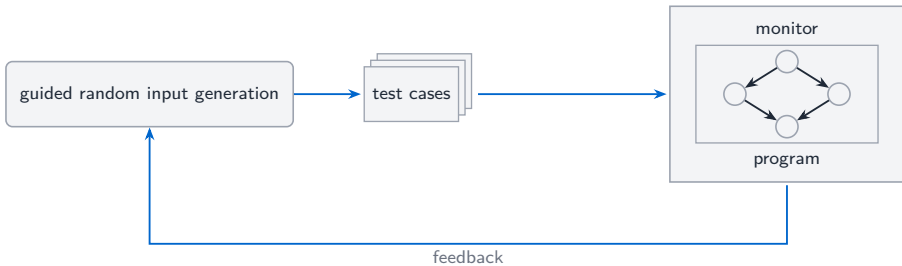


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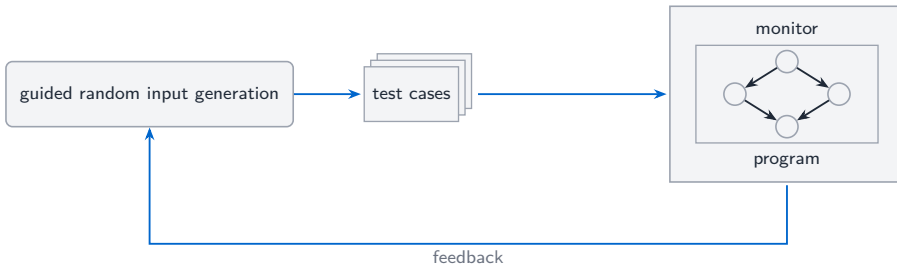


The generator sees nothing of the program; the monitor reports only whether a run terminated abnormally. Inputs are produced independently of every run that preceded them.

# Grey-box Fuzzing



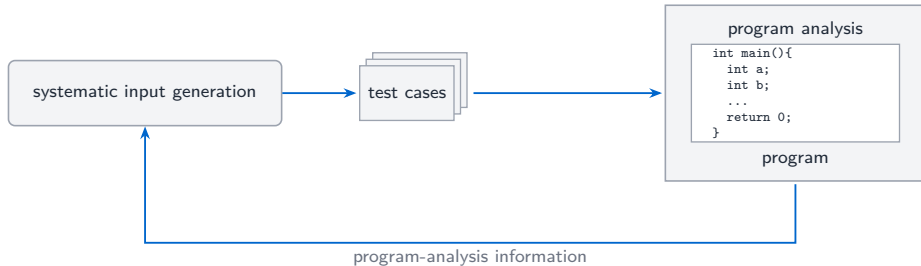
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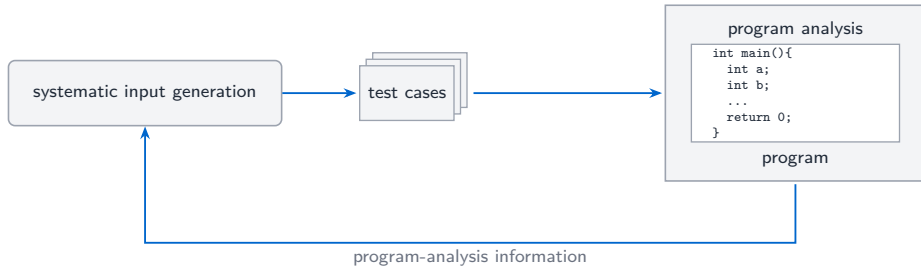
The monitor also reports which parts of the program the run executed, and the generator uses that report when choosing the next input.



# White-box Fuzzing

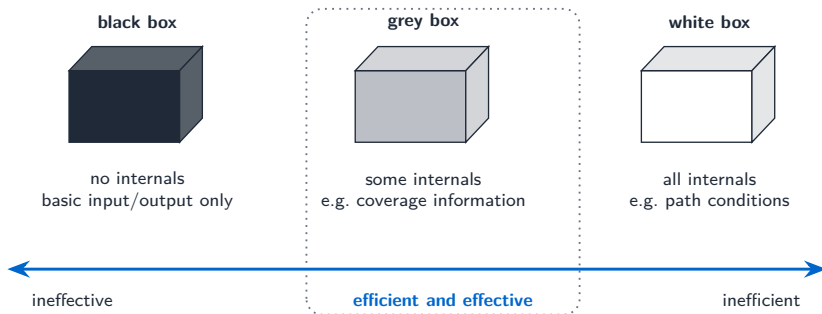


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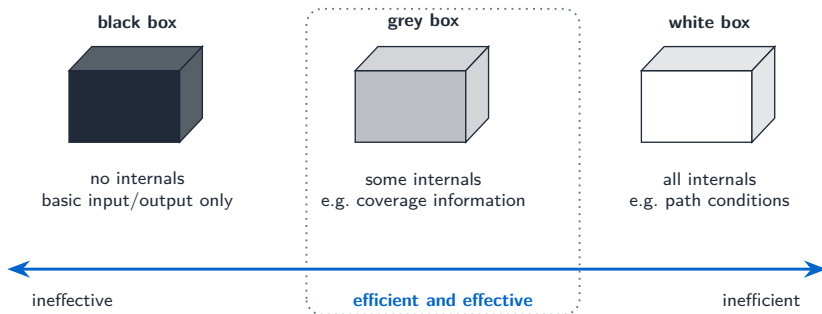


The generator reasons about the program text itself — path conditions, symbolic values — and derives inputs that reach a chosen branch.

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Most fuzzers in use combine elements of all three.

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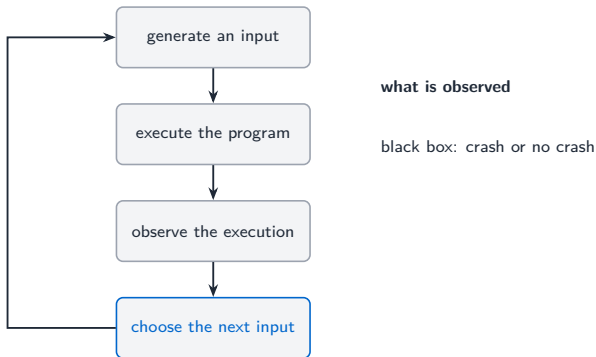
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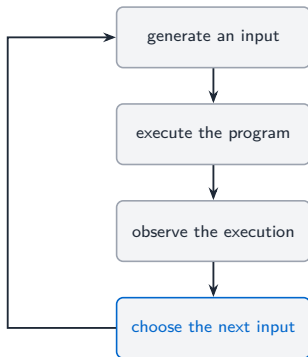
Constraint solving is expensive and does not scale to every path—>grey-box techniques.



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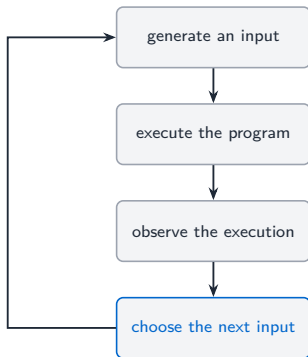


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**what is observed**

black box: crash or no crash

grey box: new coverage or not

also: a new path, a new state,  
progress on a comparison