## AutoSense : A Framework for Automated Sensitivity Analysis of Program Data

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11th European Symposium on Software Engineering and Foundations of Software Engineering (ESEC-FSE), Paderborn, Germany. September 6, 2017

### OUTLINE

Sensitivity Analysis: A Motivating Example

Problem Statement and Contributions

Definition: Sensitive Data

Dynamic Sensitivity Analysis

Static-Dynamic Combined Sensitivity Analysis

Results Reliability of Sensitivity Analysis Evaluation of Dynamic Analysis Evaluation of Static-Dynamic Combined Analysis

Conclusion

## Sensitivity Analysis: A Motivating Example

```
bool binsearch (int lo, int hi)
ł
    unsigned int size = hi - lo + 1;
    unsigned int mid = (lo+hi)/2;
    if (lo>hi) return false;
    if (size \geq 1)
        if(a[mid] == key) return true;
        else if(a[mid]>key)
              return binsearch (lo, mid -1);
        else return binsearch (mid+1, hi);
    return false;
}
   Sensitive
               Insensitive
```

Note that any erroneous value except 0 that size may take, the program behavior remains unaffected.

### Problem Statement and Contributions



Figure 1: Data classification in approximate computing

#### **Contributions:**

- A Dynamic analysis to automatically classify program data as sensitive or insensitive with probabilistic guarantee.
- A Static-Dynamic combined analysis for efficiency.

#### Definition: Sensitive Data

Given an acceptable QoS band for a program  $\mathcal{P}$  and a sensitivity threshold probability  $\theta$ , a program data  $v \in \mathcal{D}$  is called sensitive if and only if  $\forall e \in E$ , the probability that the program output remains in the acceptable QoS band when every instance  $(v_e, \ell)$  in e is replaced with some  $(v_{approx}, \ell)$ , is less than  $\theta$ .

$$SD = \{ v \in D \mid \forall e \in E, \forall \ell \in \ell_v^e, (v_e, \ell) \to (v_{approx}, \ell) \implies Pr(\mathcal{R} \in QoS) < \theta \}$$
(1)

where  $(v_e, \ell) \rightarrow (v_{approx}, \ell)$  means the substitution of  $(v_{approx}, \ell)$  in place of  $(v_e, \ell)$ . The set of *insensitive* data is defined as  $\overline{SD} = D - SD$ .



### Dynamic Sensitivity Analysis

For every  $v \in \mathcal{D}$ , we propose a hypothesis that  $\forall e \in E, \forall \ell \in \ell_v^e, (v_e, \ell) \rightarrow (v_{approx}, \ell) \implies \mathcal{R} \in QoS$ . Let us denote such an hypothesis by K. Test the following null and contrary hypothesis:

$$H: Pr(K) < \theta; H': Pr(K) \ge \theta$$
(2)

where Pr(K) is the probability that the hypothesis K is true.



Figure 2: Framework of Dynamic Sensitivity Analysis with Hypothesis Testing

# Static-Dynamic combined Sensitivity Analysis

#### Limitation of Dynamic Sensitivity Analysis

- Compute and data intensive programs may take a long time to terminate, making each trial during the hypothesis testing expensive.
- Generating random inputs for many applications can be challenging.

#### Static-Dynamic combined analysis

- Based on classical data flow analysis.
- The sensitivity of some variables like global data, method parameters and method local data whose expression has constant(s) or function call(s) is initialized using dynamic analysis.
- The principle behind our static analysis is that insensitive data cannot flow into sensitive data.

## Static-Dynamic Combined Sensitivity Analysis

#### Definition: Data Sensitivity Lattice

The elements of the complete lattice *L* of our analysis are mappings  $\sigma : \mathcal{D} \to \{\bot, S, I, \top\}.$ 

- $\sigma(x) = \bot$  denotes that no information is known about the data x
- $\sigma(x) = \top$  denotes that x may be *sensitive* or *insensitive*.
- $\sigma(x) = S$  and  $\sigma(x) = I$  denotes x to be sensitive and insensitive respectively.

We define a *data sensitivity lattice* over the range of  $\sigma$ , i.e.,  $\{\bot, S, I, \top\}$ 



Figure 3: Data Sensitivity Lattice

### Static-Dynamic Combined Sensitivity Analysis

#### Partial order on $\sigma$

The partial order on  $\sigma$  is defined :

$$\forall \sigma : \bot_{\sigma} \sqsubseteq \sigma \forall \sigma_1, \sigma_2 : \sigma_1 \sqsubseteq \sigma_2 \text{ iff } \forall x, \sigma_1(x) \sqsubseteq_D \sigma_2(x).$$

$$(3)$$

where  $\perp_{\sigma} \in \sigma$  maps every  $x \in D$  to  $\perp$ ,  $\sqsubseteq$  denotes the partial order relation on  $\sigma$  and  $\sqsubseteq_D$  denotes the partial order relation of the *data* sensitivity lattice.

The *join* operation over  $\sigma$  is defined in Eq. 4.

$$(\sigma_1 \sqcup \sigma_2)(x) = \sigma_1(x) \sqcup_D \sigma_2(x) \tag{4}$$



Figure 4: Data Sensitivity Lattice

#### Transfer Functions

Considering a general assignment statement block [x := a], *a* being any expression, we define the transfer functions of our analysis as:

$$[x = a]: f(\sigma) = \begin{cases} \sigma(x \to I) & \text{if } \forall v \in FV(a), \sigma(v) = I \\ \sigma(x \to S) & \text{if } \forall v \in FV(a), \sigma(v) = S \\ \sigma(x \to T) & \text{if } \exists u, v \in FV(a) \\ & \text{s.t. } \sigma(u) = S, \sigma(v) = I \\ \sigma & \text{if } FV(a) = \emptyset \end{cases}$$
(5)  
$$[\cdots]: f(\sigma) = \sigma$$

where  $[\cdots]$  is to denote any program statement which is not an assignment statement and FV(a) is the set of all free variables of the expression *a*.

## Reliability of Sensitivity Analysis



Figure 5: Percent output failing QoS with confidence  $\theta = 0.3$  and  $\theta = 0.5$ 



Figure 6: Raytracer rendered image with AutoSense guided approximation

### Evaluation of Dynamic Analysis



Figure 7: The percent insensitive data reported by a AutoSense on varying QoS  $\gamma$  and fixed probability factor  $\theta=0.5$ 

### Evaluation of Dynamic Analysis



Figure 8: The percent insensitive data reported by a AutoSense on varying  $\theta$  and fixed QoS  $\gamma = 0.5$  (scimark2), PSNR=10.5 (raytracer ) and exact (jmeint )

### Evaluation of Static-Dynamic Combined Analysis



Figure 9: Performance of Static-dynamic combined vs. Dynamic analysis

Application	TP	FP	FN	Precision (%)	Recall (%)
FFT	0	0	3	0	0
SOR	3	0	0	100	100
MC	1	0	1	100	50
SMM	2	0	0	100	100
LU	0	0	9	0	0
Raytracer	0	1	2	0	0

Table 1: Precision, Recall of the Combined Analysis w.r.t. Dynamic Analysis

- Identification of insensitive error resilient data of an application is non-trivial, especially when the application is large and has substantial data and control dependencies.
- We Illustrate that a systematic study of the effect of inaccuracy in program data with statistical methods like hypothesis testing can lead to automatic classification of insensitive and sensitive data.
- Dynamic analysis is computationally expensive and time consuming for many applications. We propose a static analysis to derive insensitive data, with efficiency. Although static analysis shows high precision w.r.t. dynamic analysis, it fails to identify many.

#### THANK YOU ...